

HW2 - Formal grammars (Solution)

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1. What are the languages engendered by the following context-free grammars:

(a) [1.5 pts.] $\mathcal{G}_1 = (\{a, b\}, \{S\}, S, \{S \rightarrow aSa \mid bSb \mid \varepsilon\})$

Solution: $L(\mathcal{G}_1) = \{ww^R \mid w \in \{a, b\}^*\}$.

(b) [2.5 pts.] $\mathcal{G}_2 = (\{a, b\}, \{S, T\}, S, \{S \rightarrow aSa \mid bSb \mid aTb \mid bTa, T \rightarrow Ta \mid Tb \mid \varepsilon\})$

Solution: $L(\mathcal{G}_2) = \{w \in \{a, b\}^* \mid w \neq w^R\}$.

2. Give a context-free grammar for each of the following languages:

(a) [2.5 pts.] $\{a^m b^n c^p \mid n \geq m + p\}$

Solution:

$$\begin{aligned} S &\rightarrow ABC \\ A &\rightarrow aAb \mid \varepsilon \\ B &\rightarrow bB \mid \varepsilon \\ C &\rightarrow bCc \mid \varepsilon \end{aligned}$$

(b) [1 pts.] $\{va^n \mid v \in \{a, b\}^n \wedge |v|_a = n\}$

Solution: $S \rightarrow aSa \mid \varepsilon$

There was an error in the question; the intended language was

$$\{va^n \mid v \in \{a, b\}^* \wedge |v|_a = n\}$$

$$\begin{aligned} S &\rightarrow BaSa \mid B \\ B &\rightarrow bB \mid \varepsilon \end{aligned}$$

For those who tried (and almost got it right), this question awarded up to 2.5 points.

(c) [2.5 pts.] $\{w \in \{a, b\}^* \mid |w|_a \geq |w|_b\}$

Solution: $S \rightarrow aSbS \mid bSaS \mid aS \mid \varepsilon$

3. [5 pts.] Propose a context-free grammar that generates the set of valid arithmetic expressions (in $\mathbb{N}$) over $\{0, 1, \dots, 9, \times, +, (,)\}$ with the two following restrictions:

- the numbers should not contain any leading 0's, except for 0 itself, and
- there should not be any redundant parentheses.

For instance, expressions in (1) are valid, while those in (2) are not.

- (1)
- a. $1234 \times 0 \times (1 + 654 + 292929)$
 - b. $4 \times 911 + 2 \times (42 + 5555 \times 886) \times 1$
 - c. $9 \times (1 + 2 + 7) \times (4 \times 2 + 1) + 5 \times 112$
 - d. $(0 + 4) \times (5 \times 7 + 91) + 4$

- (2)
- a. $*010$
 - b. $*1 \times (2)$
 - c. $*4 \times 911 + (2 \times 42 + 5555)$
 - d. $*9 \times (1 \times 2) + 74$

Is your grammar ambiguous? If so, show two parse trees for a valid expression; otherwise, explain informally why your grammar is not (a formal proof is not expected).

Solution:

$$\begin{aligned} E &\rightarrow E + M \mid M \\ M &\rightarrow M \times N \mid (A) \times M \mid M \times (A) \mid (A) \times (A) \mid N \\ A &\rightarrow E + M \\ N &\rightarrow 0 \mid 1N' \mid 2N' \mid \dots \mid 9N' \\ N' &\rightarrow 0N' \mid 1N' \mid 2N' \mid \dots \mid 9N' \mid \varepsilon \end{aligned}$$

This grammar is ambiguous due to the production rules for M . For instance, the expression $(1 + 2) \times (2 + 0) \times 1$ has two derivations according to when ‘ $\times 1$ ’ is derived.

4. [5 pts.] Define a non-trivial context-free grammar that could generate the sentences in (3).
- (3)
- Alice eats a sweet.
 - The waiter gives Alice cakes in a tray.
 - The tabby cat with a grin disappears.
 - Alice likes that her brother plays football with his friends.

Solution:

S	→	NP VP
NP	→	PropN
		Det N'
		N'
N'	→	Adj N'
		N' PP
		N
PP	→	P NP
VP	→	VP PP
		V _t NP
		V _d NP NP
		V _t <i>that</i> S
		V _i
PropN	→	<i>Alice</i>
Det	→	<i>a the her his</i>
N	→	<i>sweet waiter cakes tray cat grin brother football friends</i>
Adj	→	<i>tabby</i>
P	→	<i>in with</i>
V _t	→	<i>eats likes plays</i>
V _d	→	<i>gives</i>
V _i	→	<i>disappears</i>

Note that the rules $N' \rightarrow \text{Adj } N' \mid N' \text{ PP}$ allow for the ambiguity in (3c):

- Either there is only one *tabby cat* with a grin: $N' \rightarrow \text{Adj } N' \rightarrow \text{Adj } N' \text{ PP}$, or
- there is only one *cat with a grin* that is tabby: $N' \rightarrow N' \text{ PP} \rightarrow \text{Adj } N' \text{ PP}$